

# The Halting Problem

Joseph Paul Cohen

## What is the halting problem?

Given some program and some input, determine if the program will halt on the given input.

$HALT(\#_p, x)$  = Does the program  $\#_p$  halt on input  $x$

$$HALT(\#_p, x) = \begin{cases} TRUE & \text{When } \#_p \text{ is run on } x \text{ it ends} \\ FALSE & \text{When } \#_p \text{ is run on } x \text{ it does not end} \end{cases}$$

## What is a program?

Some set of instructions that operate on some input

```
1  $f(x)$  :  
2 while  $x < 3$  do  
3   | print  $x$   
4   |  $x = x + 1$ 
```

```
1  $f(x)$  :  
2 while  $x < 3$  do  
3   | print  $x$   
4   |  $x = x + 1$ 
```

$$f(0) \longrightarrow 0, 1, 2$$

$$f(-1) \longrightarrow -1, 0, 1, 2$$

$$f(5) \longrightarrow \text{NOTHING PRINTED}$$

What is a program?

Every program can be represented as some number

$f(x)$  can be represented as some number  $\#_f$ , such as 239847238

$$f(x) \longrightarrow \#_f$$

```
1  $f(x)$  :  
2 while  $x < 3$  do  
3   |   print  $x$   
4   |    $x = x + 1$ 
```

What does this return?  
 $HALT(\#_f, 2)$

```
1  $f(x)$  :  
2 while  $x < 3$  do  
3   | print  $x$   
4   |  $x = x + 1$ 
```

What does this return?

$HALT(\#_f, 2)$

TRUE

```
1  $t(x)$  :  
2 while  $x > 3$  do  
3   |   print  $x$   
4   |    $x = x + 1$ 
```

What does this return?  
 $HALT(\#_t, 2)$



```
1  $t(x)$  :  
2 while  $x > 3$  do  
3   | print  $x$   
4   |  $x = x + 1$ 
```

What does this return?

$HALT(\#_t, 2)$

TRUE

```
1  $t(x)$  :  
2 while  $x > 3$  do  
3   |   print  $x$   
4   |    $x = x + 1$ 
```

What does this return?  
 $HALT(\#_t, 5)$

```
1  $t(x)$  :  
2 while  $x > 3$  do  
3   | print  $x$   
4   |  $x = x + 1$ 
```

What does this return?

$HALT(\#_t, 5)$

FALSE

There can never be a program for the the function *HALT*

There can never be a program for the the function *HALT*

This is a big deal

## Outline for Proof

We assume that a program for *HALT* exists. (There is a  $\#_{HALT}$ )

We write a program with  $\#_{HALT}$

We use that program and come to a contradiction

$$g(x) = \begin{cases} \text{FALSE} & \text{If } x \text{ run on } x \text{ never halts} \\ \uparrow & \text{Never returns if } x \text{ run on } x \text{ halts} \end{cases}$$

---

```

1 g(x) :
2 if HALT(x,x) == TRUE then
3   | while TRUE do
4   |   | print "Never Returns"
5 else
6   | return FALSE

```

---

$$g(x) = \begin{cases} FALSE & \text{If } x \text{ run on } x \text{ never halts} \\ \uparrow & \text{Never returns if } x \text{ run on } x \text{ halts} \end{cases}$$

```

1 g(x) :
2 if HALT(x,x) == TRUE then
3   | while TRUE do
4   |   | print "Never Returns"
5 else
6   | return FALSE

```

$$HALT(\#_p, x) = \begin{cases} TRUE & \text{If when } \#_p \text{ is run on } x \text{ it ends} \\ FALSE & \text{If when } \#_p \text{ is run on } x \text{ it does not end} \end{cases}$$



$g$  is a program so it has a number  $\#_g$

What happens when we run  $\#_g$  on itself?

$g(\#_g)$

```
1  $g(x)$  :  
2 if  $HALT(x,x) == TRUE$  then  
3   |   while  $TRUE$  do  
4   |   |   print “Never Returns”  
5 else  
6   |   return  $FALSE$ 
```

Lets look at what happens when:  $g$  halts when run on  $g$

$\implies HALT(\#_g, \#_g) == FALSE$

$\implies g$  does not halt on  $g$

```
1  $g(x)$  :  
2 if  $HALT(x,x) == TRUE$  then  
3   | while  $TRUE$  do  
4   |   | print “Never Returns”  
5 else  
6   | return  $FALSE$ 
```

How about:  $g$  does not halt when run on  $g$

$\implies HALT(\#_g, \#_g) == TRUE$

$\implies g$  halts when run on  $g$

We have a contradiction so there can be no program for *HALT*